

About Machine Readable Travel Documents

Privacy Enhancement Using (Weakly) Non-Transferable Data Authentication

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LASEC

1 ICAO-MRTD

2 Non-Transferable Proofs

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1 ICAO-MRTD

- ICAO-MRTD Overview
- Data Structures and PKI
- MRTD Cryptography
- Security and Privacy Issues

2 Non-Transferable Proofs

Objectives

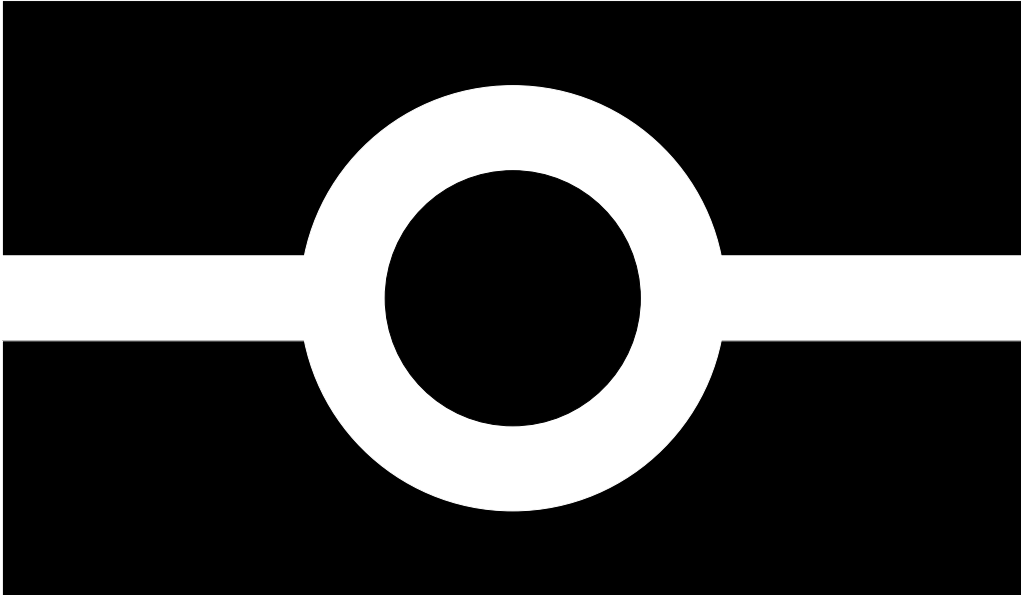
to enable inspecting authorities to securely identify visitors with the help of machine-readable digital information

- biometrics
 - contactless IC chip
 - digital signature + PKI
-
- maintained by UN/ICAO (International Civil Aviation Organization)

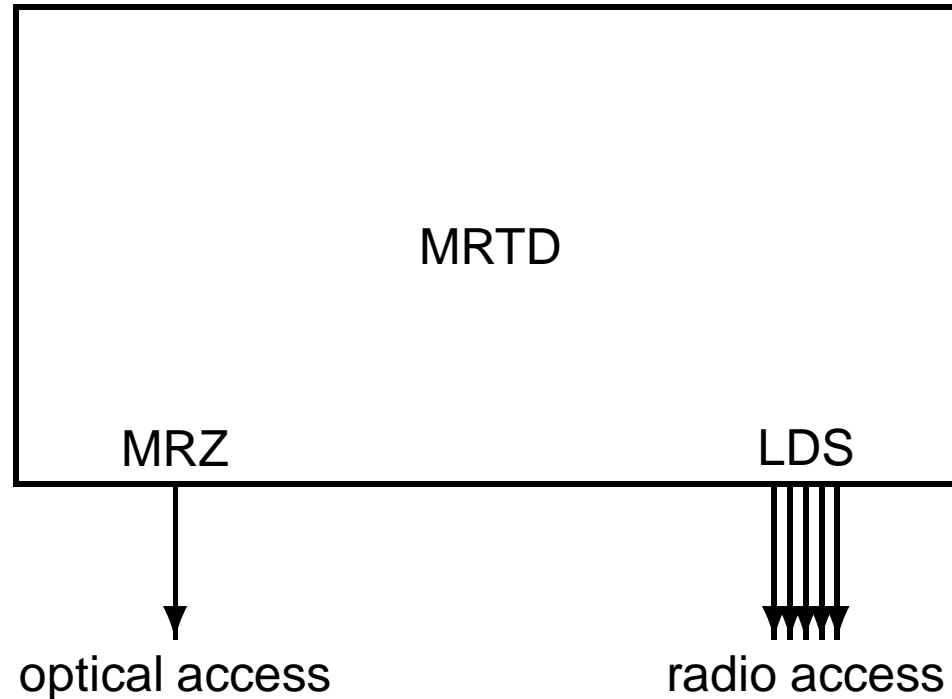
MRTD History

- 1968: ICAO starts working on MRTD
- 1980: first standard (OCR-B Machine Readable Zone (MRZ))
- 1997: ICAO-NTWG (New Tech. WG) starts working on biometrics
- 2001 9/11: US want to speed up the process
- 2004: version 1.1 of standard with ICC
- 2006: extended access control under development in the EU

How to Distinguish a Compliant MRTD



MRTD in a Nutshell



- data authentication by digital signature + PKI
aka **passive authentication**
- access control + key agreement based on MRZ_info
aka **basic access control (BAC)**
- chip authentication by public-key cryptography
aka **active authentication (AA)**

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MRZ Example

- document type
- issuing country
- holder name
- doc. number + CRC
- nationality
- date of birth + CRC
- gender
- date of expiry + CRC
- options + CRC

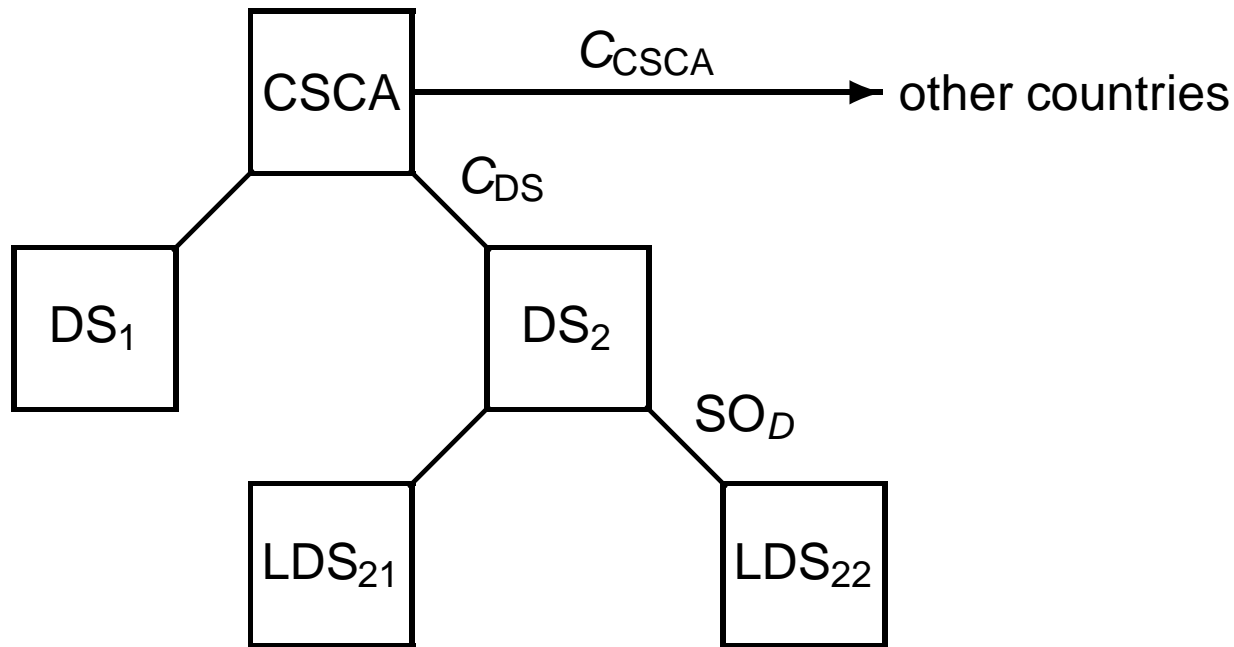
LDS Structure

- **DG1 (mandatory): same as MRZ**
- **DG2 (mandatory): encoded face**
- DG3: encoded finger(s)
- DG4: encoded eye(s)
- DG5: displayed portrait
- DG6: (reserved)
- DG7: displayed signature
- DG8: data feature(s)
- DG9: structure feature(s)
- DG10: substance feature(s)
- DG11: add. personal detail(s)
- DG12: add. document detail(s)
- DG13: optional detail(s)
- DG14: (reserved)
- DG15: $KP_{u_{AA}}$
- *DG16: person(s) to notify*
- DG17: autom. border clearance
- DG18: electronic visa
- DG19: travel record(s)
- **SO_D (mandatory)**

SO_D Structure

- list of hash for data groups DG1–DG15
- formatted signature by DS (include: information about DS)
- (optional) C_{DS}

Hierarchy



- one PKI per country
one CSCA (Country Signing Certificate Authority)
 C_{CSCA} : self-signed CSCA public key $K_{Pu_{CSCA}}$
 C_{CSCA} distributed to other countries by diplomatic means
- possibly many DS (Document Signer) per country
 C_{DS} : certificate for a DS public key $K_{Pu_{DS}}$
- SO_D : signature of (part of) LDS in MRTD

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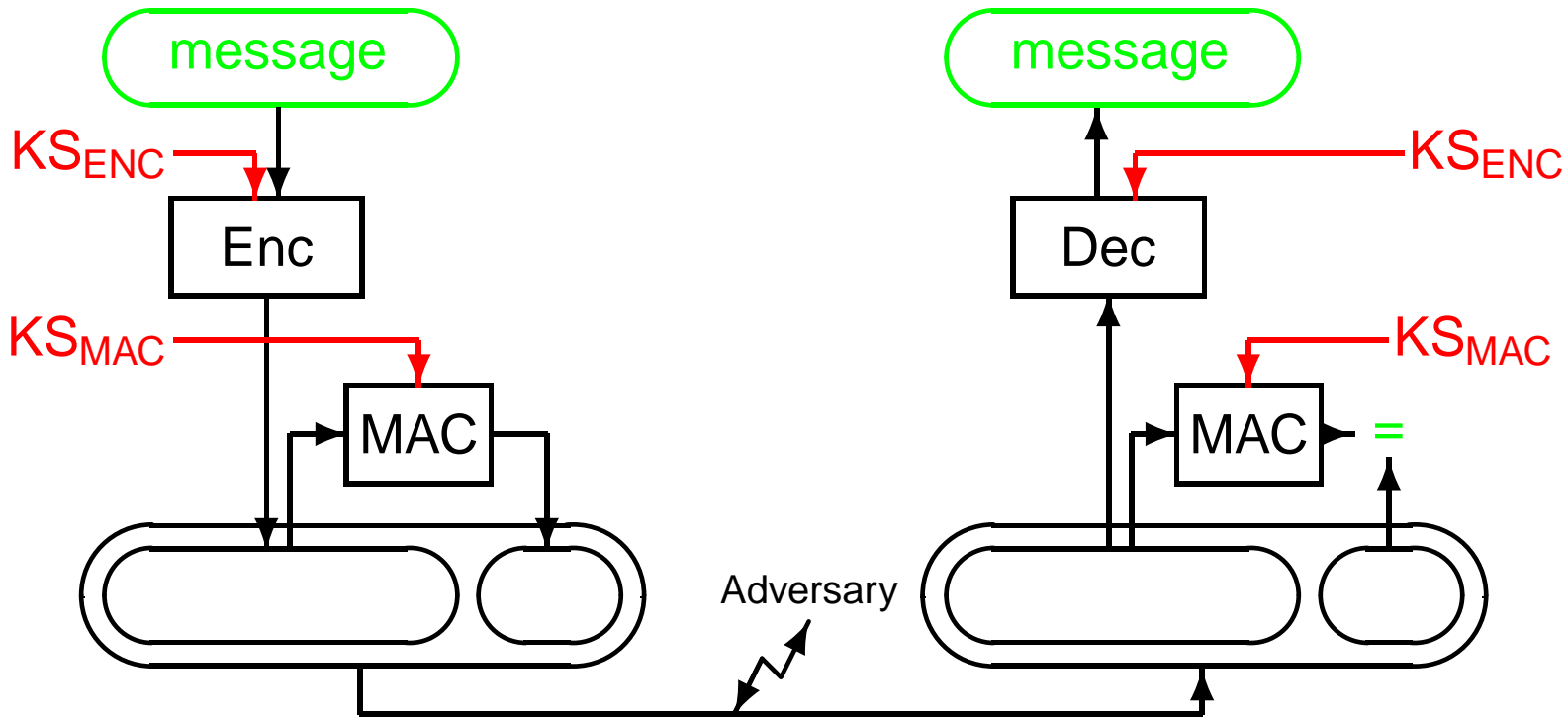
2 Non-Transferable Proofs

Basic Access Control

- goal** prevent from unauthorized access by the holder (privacy)
- read MRZ (OCR-B)
 - extract MRZ_info
 - run an authenticated key exchange based on MRZ_info
 - open secure messaging based on the exchanged symmetric key
- proves that reader knows MRZ_info

Secure Messaging

goal authentication, integrity, confidentiality of communication



→ secure channel based on 3DES

Passive Authentication

goal authenticate LDS

- after getting SO_D , check the included certificate C_{DS} and the signature
 - when loading a data group from LDS, check its hash with what is in SO_D
- stamp by DS on LDS

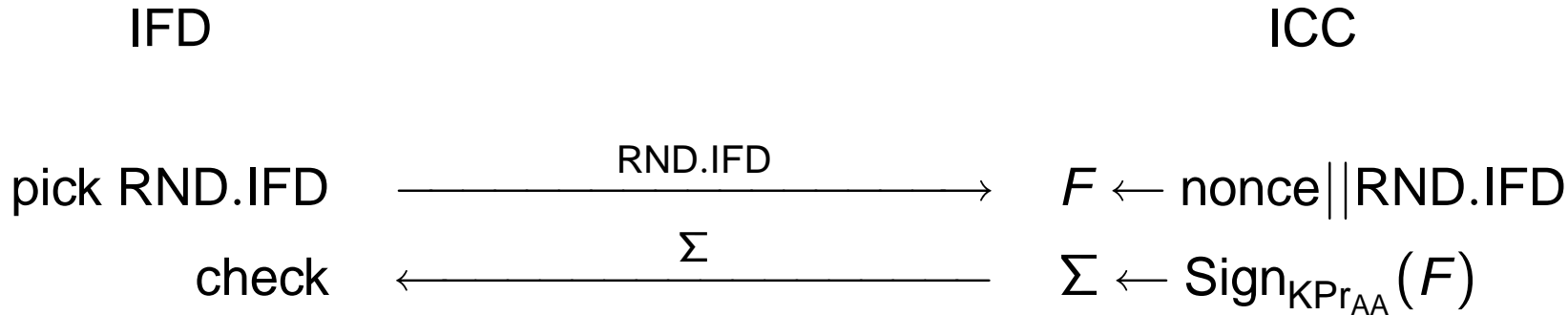
Active Authentication

goal authenticate the chip

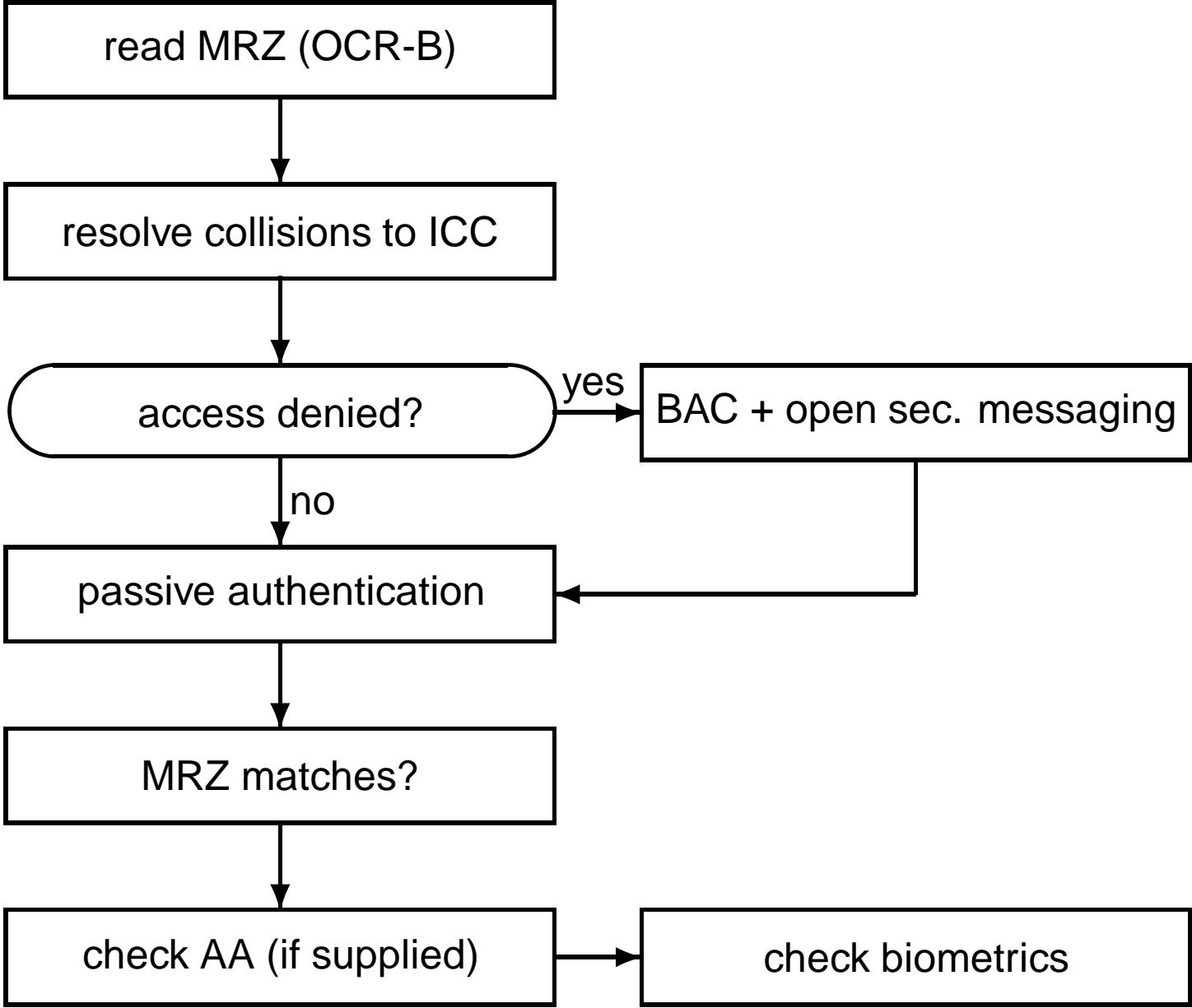
- proves that ICC knows some secret key KPr_{AA} linked to a public key KPu_{AA} by a challenge-response protocol
(KPu_{AA} in LDS authenticated by passive authentication)

→ prove that the chip is not a clone

Active Authentication Protocol



Sequence of Steps for Identification



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Coming From Wireless Technology

(claimed to be possible at a distance of 10m)

- detecting the proximity of an e-passport
 - threat:** giving valuable information to passport theafs
 - threat:** privacy (in some cases) by tracking people
- data skimming
 - threat:** privacy
- unauthorized access
 - threat:** privacy

Coming From IC Chip

- too much trust in automated process, lazy identification
threat: identity theft
- malicious cookies put in MRTD
threat: privacy
- dependence on the technology: DoS attack could kill the IC chip
threat: waste of time at border controls
- abuse of automatic recognition
threat: privacy
- **leakage of digital evidence**
threat: privacy

Digital Evidence: Challenge Semantics Attack

challenge semantics in AA:

- evidence that D existed when MRTD was queried

$$\text{RND.IFD} = H(D)$$

$$\text{evidence} = (D || \text{LDS} || \Sigma)$$

- evidence that MRTD was accessed at time t

$$\text{RND.IFD} = H(\text{social}(t - 1))$$

$$\text{evidence} = \text{timestamp}_t(\text{social}(t - 1) || \text{LDS} || \Sigma)$$

Digital Evidence: Transferable LDS Authentication

- signed personal data (name, age, gender, face, etc)
- can no longer hide/deny name, age, gender...
- when DG11 is used: more personal data (place of birth etc)
- personal profiles can be sold if they come with a proof

1 ICAO-MRTD

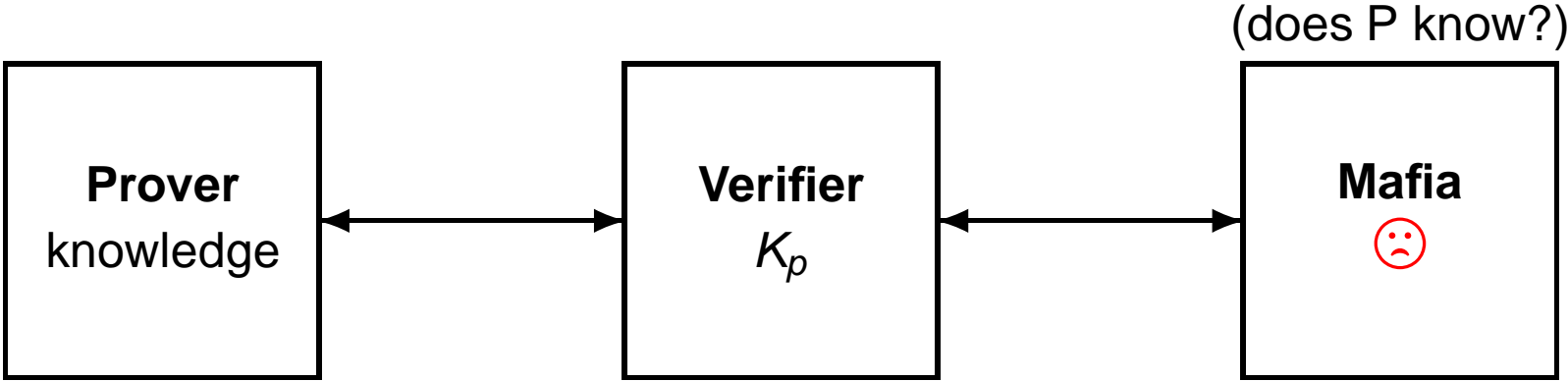
2 Non-Transferable Proofs

1 ICAO-MRTD

2 **Non-Transferable Proofs**

- Notions of Non-Transferability
- ZK Protocols for MRTD

Mafia Fraud + Fully Non-Transferable Proof



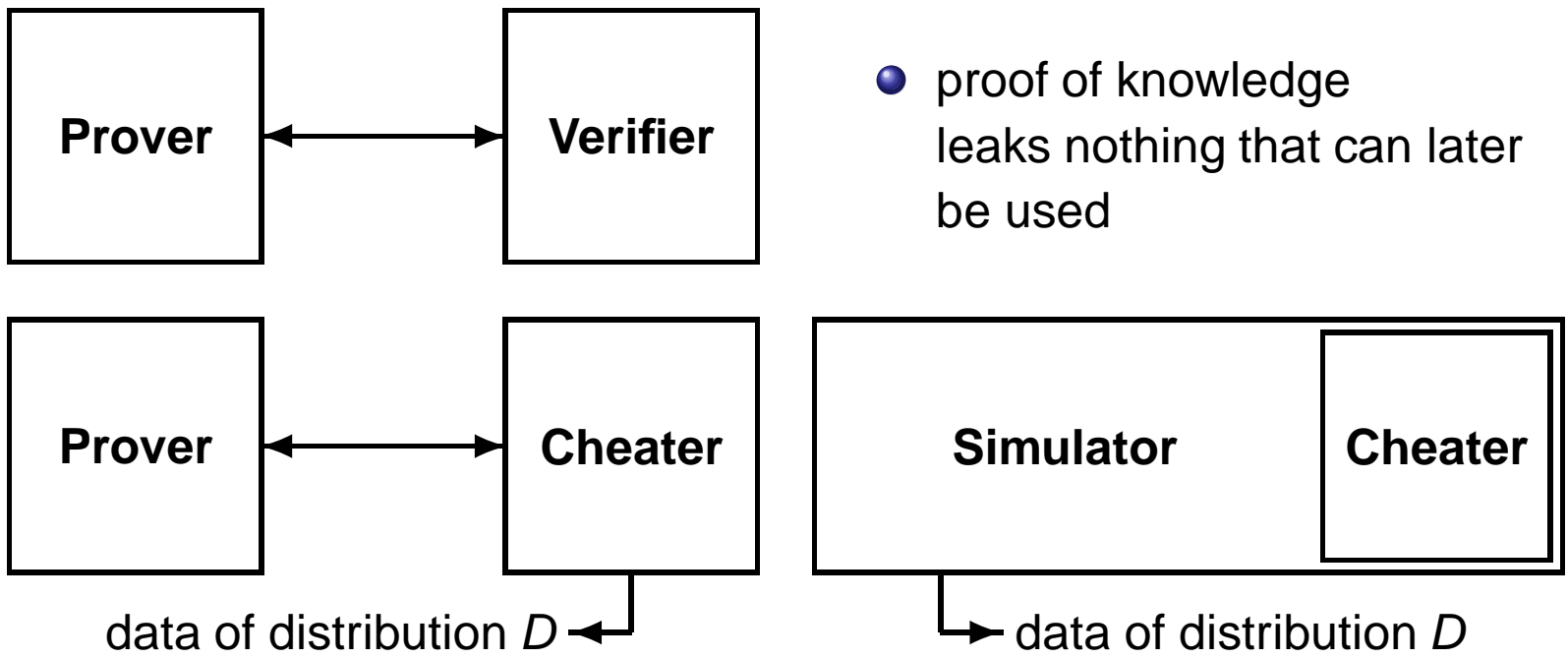
proof of knowledge



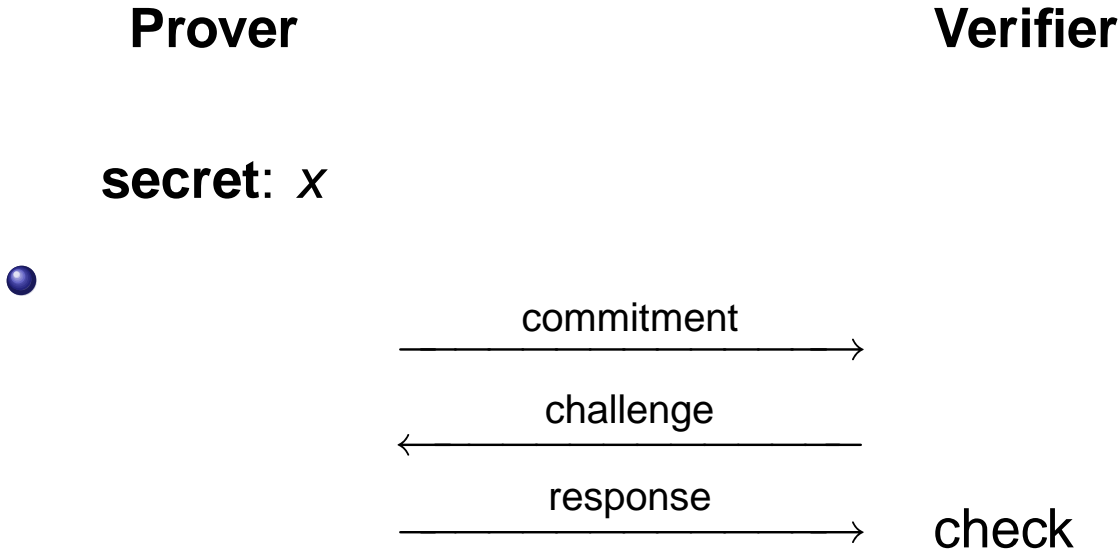
proof of knowledge or of knowing a secret key attached to K_p

→ need PKI for verifiers: maybe an overkill

Zero-Knowledge: Offline Non-Transferability



Sigma Protocols



Example: GPS Identification

Prover

parameters: g, A, B, S

public key: I ($I = g^S$)

secret key: $s \in [0, S]$

pick $r \in [0, A - 1]$

$x \leftarrow g^r$

$y \leftarrow r + cs$

\xrightarrow{x}

\xleftarrow{c}

\xrightarrow{y}

Verifier

input: I, g, A, B, S

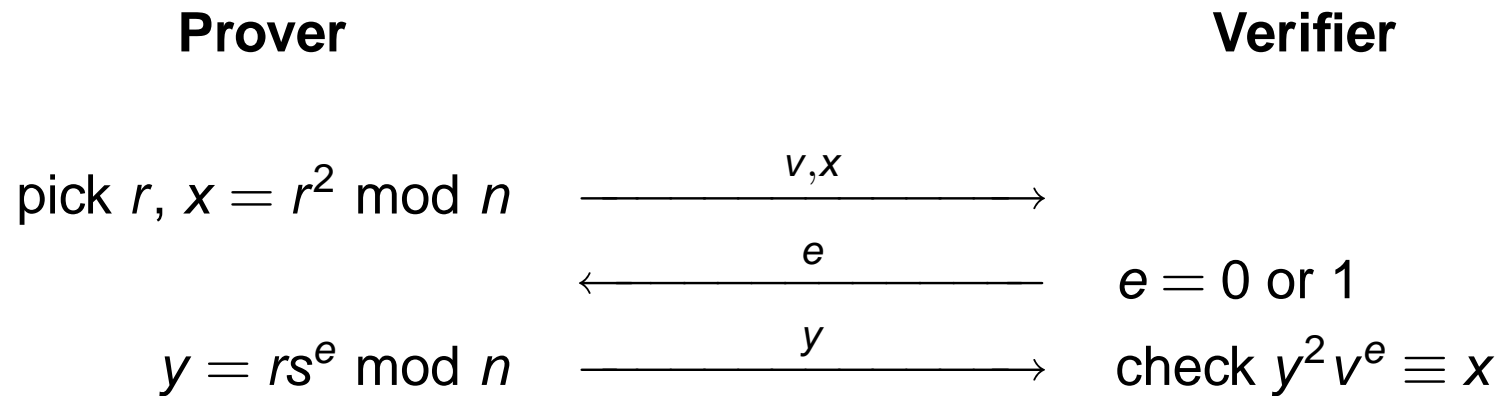
pick $c \in [0, B - 1]$

check $g^y = xI^c$

and $0 \leq y < A + (B - 1)(S - 1)$

Fiat-Shamir Signature

Basic Fiat-Shamir identification protocol:



Conversion into a signature:

- use random coins from $H(\text{message, previously seen transcript})$
- simulate the verifier using these coins
- the signature is the final transcript

Honest vs Malicious Verifier

- for Sigma-protocols: the signature is unforgeable
- malicious verifier that simulates the previous conversion: it produces a signature
- consequences: Sigma-protocols are not ZK
- maybe honest-verifier ZK
- verifiers playing the challenge semantics are not honest
- challenge semantics in GPS identification: $c = H(\text{semantics}, x)$
- UDVSP [Baek et al. Asiacrypt05]: same

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Converting Sigma Protocols into ZK Protocols

Prover

Verifier

secret: x

pick c_P

pick c_V

$\leftarrow \gamma$

$(\gamma, \delta) \leftarrow \text{commit}(c_V)$

$\xrightarrow{\text{commitment}, c_P}$

$\leftarrow \delta, c_V$

$\text{check}(c_V, \gamma, \delta)$

$\xrightarrow{\text{response}}$

check

(challenge is $c_P \oplus c_V$)

Proof of Signature Knowledge based on GQ

Prover

formatted digest: X
signature: x

pick $y \in \mathbf{Z}_N^*$

pick $c_P \in \{0, 1\}^\ell$

$Y \leftarrow y^e \pmod N$

check(c_V, γ, δ)

$z \leftarrow yx^c \pmod N$

public key: N, e

$\longleftarrow \gamma$

$\xrightarrow{Y, c_P}$

$\longleftarrow \delta, c_V$

\xrightarrow{z}

$(c = c_P \oplus c_V)$

Verifier

formatted digest: X

pick $c_V \in \{0, 1\}^\ell$

$(\gamma, \delta) \leftarrow \text{commit}(c_V)$

check $z^e = YX^c \pmod N$

Easy AA from Previous Passive Authentication

proof of holding a signature of SOD



proof of holding a secret signature of SOD

AA based on GPS

Prover

parameters: g, A, B, S
public key: I ($I = g^S$)
secret key: $s \in [0, S]$

pick $r \in [0, A - 1]$
pick $c_P \in [0, B - 1]$
 $x \leftarrow g^r$
check(c_V, γ, δ)
 $y \leftarrow r + cs$

$\xleftarrow{\gamma}$
 $\xrightarrow{x, c_P}$
 $\xleftarrow{\delta, c_V}$
 \xrightarrow{y}

Verifier

input: I, g, A, B, S

pick $c_V \in [0, B - 1]$
 $(\gamma, \delta) \leftarrow \text{commit}(c_V)$
check $g^y = xI^c$
and $0 \leq y < A + (B - 1)(S - 1)$

$(c = c_P + c_V \text{ mod } B)$

Conclusion

- privacy threat of MRTD coming from wireless channel
- privacy threat of MRTD coming from leakage of evidence
- weakly non-transferable proofs
- proof of signature knowledge based on GQ
- fix of AA

Q & A